

BOUNDEDNESS QUESTIONS FOR CALABI-YAU'S

① GENTLE INTRO TO THE MMP

② BOUNDEDNESS

③ - ④ BOUNDEDNESS for CY'S (elliptic)
+ log CY pairs

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

ONE OF THE MAIN GOALS OF ALGEBRAIC GEOMETRY IS TO PRODUCE A COMPLETE CLASSIFICATION OF PROJECTIVE VARIETIES.

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

ONE OF THE MAIN GOALS OF ALGEBRAIC GEOMETRY IS TO PRODUCE A COMPLETE CLASSIFICATION OF PROJECTIVE VARIETIES. TO THIS END, WE CAN PARTITION VARIETIES INTO EQUIVALENCE CLASSES USING 2 POSSIBLE EQUIVALENCE RELATIONS :

① ISOMORPHISM : $X \begin{array}{c} \xrightarrow{f} \\ \xleftarrow{f^{-1}} \end{array} Y$

② BIRATIONAL ISOMORPHISM : $X \begin{array}{c} \xrightarrow{g} \\ \xleftarrow{g^{-1}} \end{array} Y$

g, g^{-1} are rational maps.

$\mathbb{C}(X) =$ field of rat'l functions on X

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

X SMOOTH PROJECTIVE VARIETY / \mathbb{C}

CANONICAL BUNDLE: $\det(\Omega'_X)$

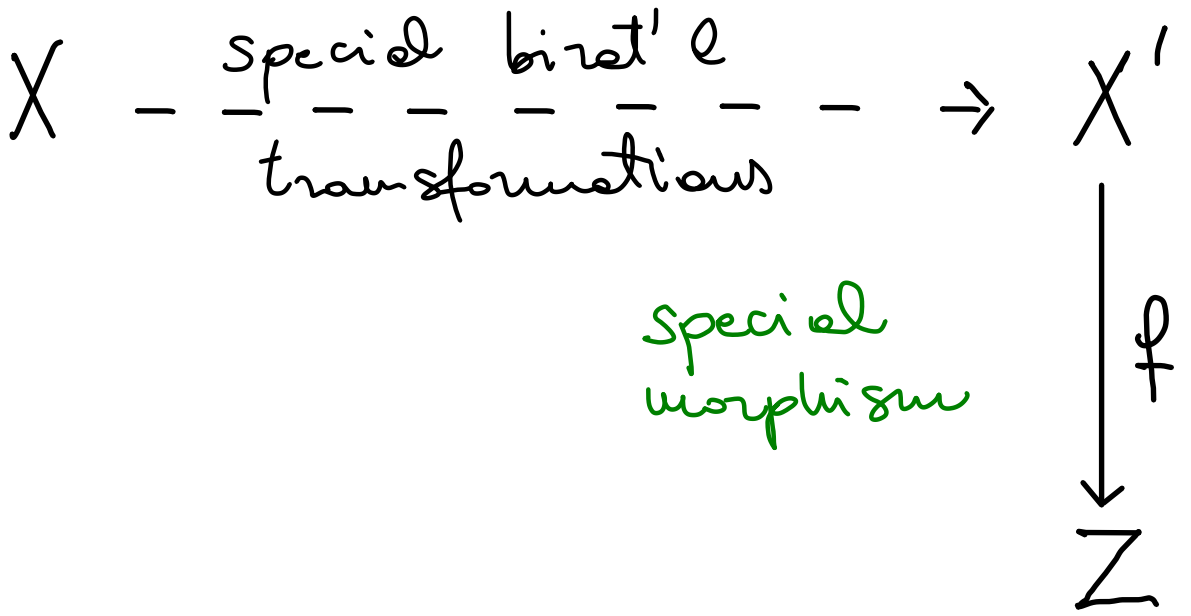
CANONICAL DIVISOR: Any Weil divisor K_X s.t.
 $\mathcal{O}_X(K_X) \cong$ canonical bundle.

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

X SMOOTH PROJECTIVE VARIETY / \mathbb{C}

CONJECTURE

STARTING FROM X , THERE EXISTS AN ALGORITHMIC WAY OF CONSTRUCTING A DIAGRAM OF THE FOLLOWING FORM



SUCH THAT f SATISFIES 1 OF THE FOLLOWING:

- ① f is BIRATIONAL & K_Z is ample Z is a can. model.
- ② f is a FIBRATION & $K_{X'} = f^* H$ ample on Z
- ③ f is a FIBR. & $-K_{X'}$ is ample on \mathbb{P}^1

the fibers will be K -trivial
 the fibers will be \mathbb{P}^1

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

HENCE IN THE BIRATIONAL CLASSIFICATION
OF ALGEBRAIC VARIETIES WE HAVE 3
IMPORTANT CLASSES OF VARIETIES THAT
ACT AS BUILDING BLOCKS

- ① CANONICAL MODELS (K is ample)
- ② K -TRIVIAL VAR'S ($K \equiv 0$)
- ③ FANO VAR'S ($-K$ is ample)

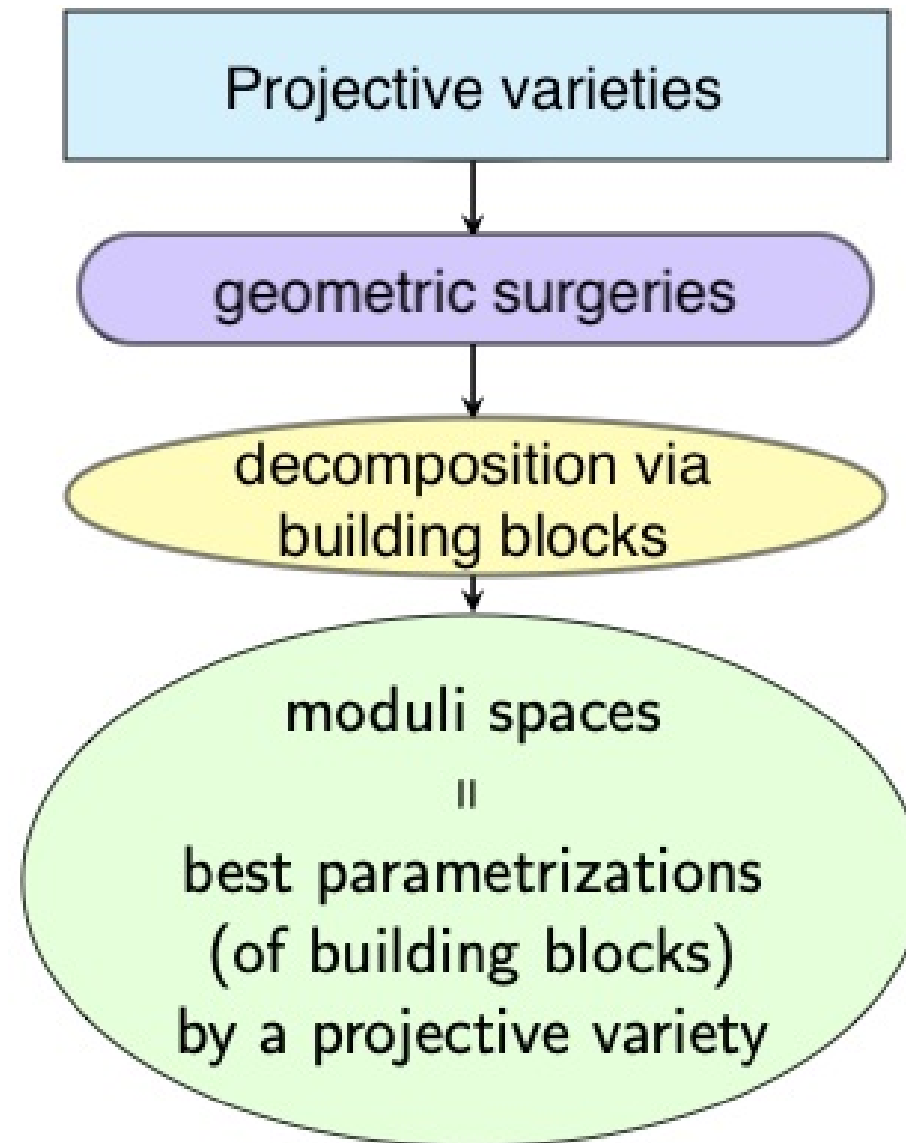
MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

THE 2nd PART OF THE CLASSIFICATION AIMS
THEN TO CLASSIFY THOSE VARIETIES THAT BELONG
TO ONE OF THE 3 BUILDING BLOCKS

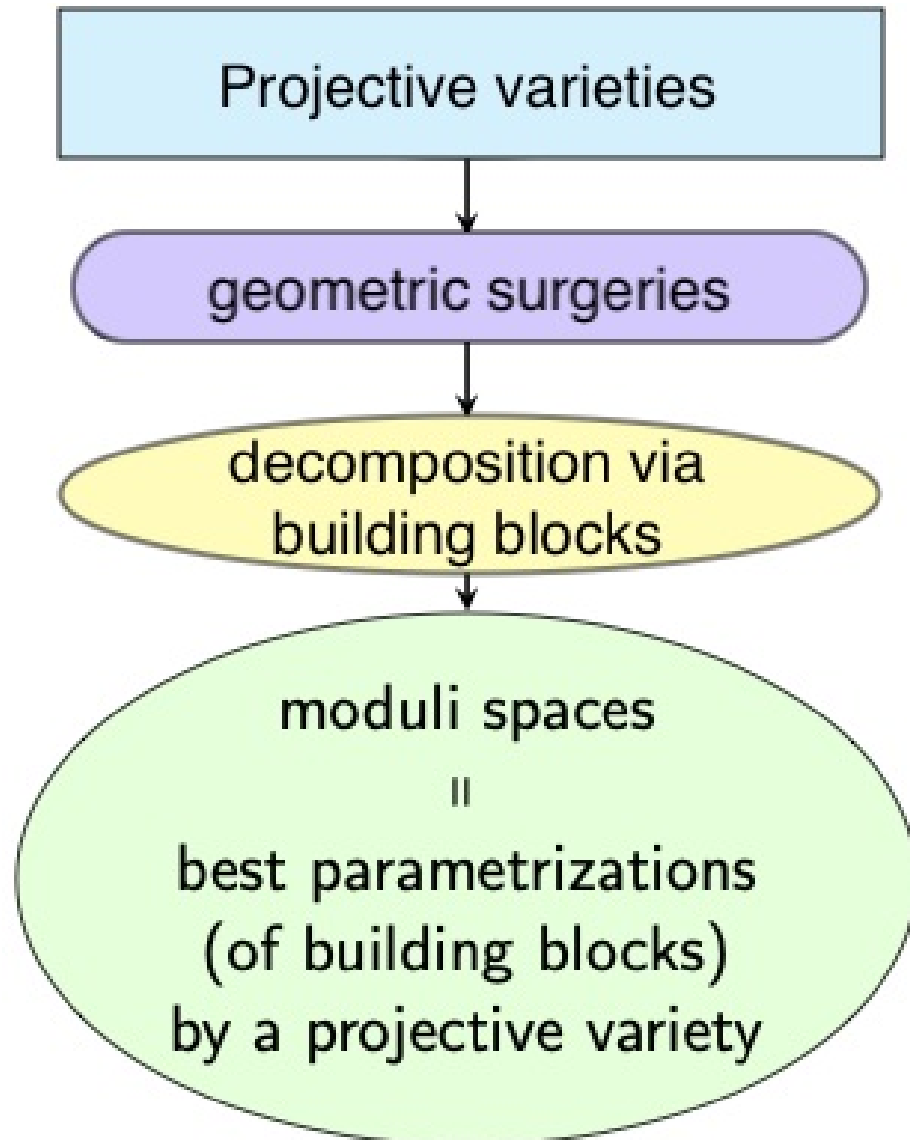
- ① CANONICAL MODELS (K IS AMPLE)
- ② K -TRIVIAL VAR'S ($K \equiv 0$)
- ③ FANO VAR'S ($-K$ IS AMPLE)

MODULI SPACES :

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION



MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION



IN ORDER TO COMPLETE SUCH CLASSIFICATION SCHEME, WE NEED ALSO TO WORK WITH SINGULAR VARIETIES.

MINIMAL MODEL PROGRAM: A GENTLE INTRODUCTION

X NORMAL PROJECTIVE VARIETY / \mathbb{C}

CANONICAL BUNDLE: $i_* \det(\Omega'_{X^{sm}})$ rank 1-sheaf

$$i: X^{sm} \hookrightarrow X$$

CANONICAL DIVISOR: Zer. closure of any

$$K_{X^{sm}}$$

K_X will be a Weil divisor
(not. nec. Cartier)

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

DEFINITION

① A PAIR IS THE DATUM OF :

(i) X a normal variety

(ii) B an effective Weil divisor with BOUNDARY
coefficients $\in (0, 1]$

② A LOG PAIR IS A PAIR (X, B) SUCH THAT

$K_X + B$ IS \mathbb{R} -CARTIER

(can be written as a \mathbb{R} -sum
of Cartier divisors)

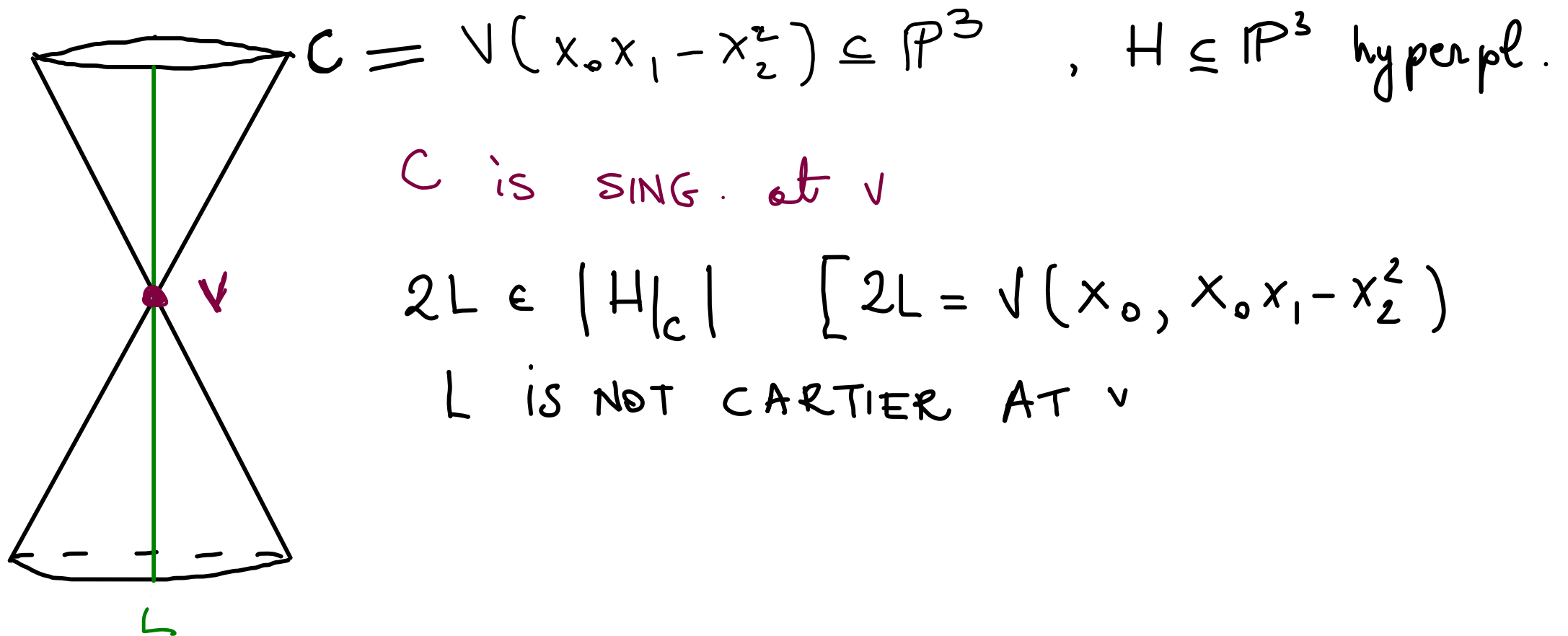
MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

WHY IS IT USEFUL TO CONSIDER LOG PAIRS?

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WHY IS IT USEFUL TO CONSIDER LOG PAIRS?

CONSIDER FOR EXAMPLE



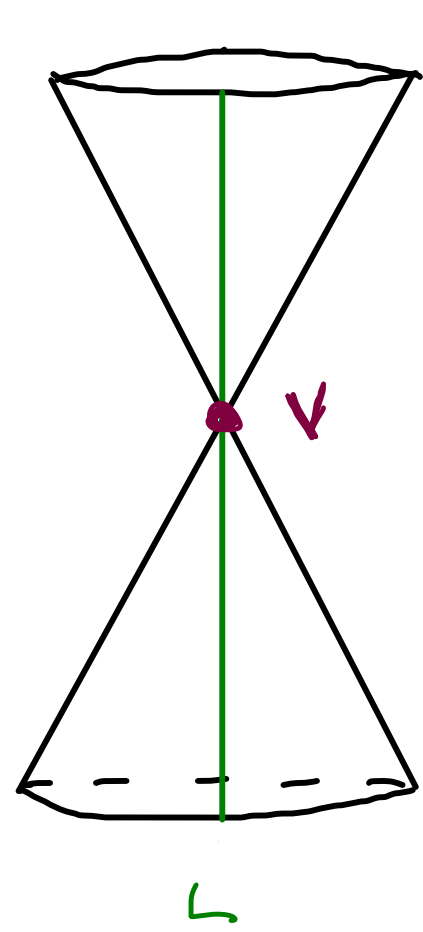
ADJ. FORM: X SMOOTH
 D^{ur} SMOOTH
codim 1

$$(K_X + D)|_D = K_D$$

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

WHY IS IT USEFUL TO CONSIDER LOG PAIRS?

CONSIDER FOR EXAMPLE



$$C = V(x_0 x_1 - x_2^2) \subseteq \mathbb{P}^3$$

$H \subseteq \mathbb{P}^3$ hyperpl.
 $2L \sim H|_C$

ADJ. FORM: X SMOOTH
 D^u SMOOTH
 codim 1

$$(K_X + D)|_D = K_D$$

$$(K_{C, (v)} + L|_{C, (v)})|_{L|_{C, (v)}} = K_{L|_{C, (v)}}$$

$$\begin{aligned} (K_C + L) \cdot L &= -\frac{3}{2} \implies (K_C + L)|_L = K_L + \frac{1}{2} v \\ \stackrel{||}{=} (-2H|_C + \frac{1}{2} H|_C) \cdot L &= \end{aligned}$$

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

WHY IS IT USEFUL TO CONSIDER LOG PAIRS?

CONSIDER A FIBRATION IN K -TRIVIAL VARIETIES

$$f : Y \longrightarrow X$$

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

WE SAW THAT WE WILL NEED TO CONSIDER SINGULAR OBJECTS.

TO DO THAT, WE NEED TO MEASURE THEIR SINGULARITIES.

WE WILL DO SO USING (BIRATIONAL) RESOLUTIONS.

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DEFINITION LET (X, B) BE A PAIR. A BIRAT'L MORPHISM

$$\pi: X' \longrightarrow X$$

IS A LOG RESOLUTION WHEN :

① π IS A PROPER RESOLUTION OF X

② $\tilde{B} + \text{Exc}(\pi)$ IS

strict transform of B on X'

SIMPLE	S
NORMAL	N
CROSSING	C

$\forall x' \in X$
locally analytically
 $(X', \tilde{B} + \text{Exc}(\pi))$
 $(\mathbb{C}^n, \sum_{i=1}^k \{z_i=0\})$

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DEFINITION LET (X, B) BE A PAIR. A BIRAT'L
MORPHISM

$$\pi : X' \longrightarrow X$$

IS A LOG RESOLUTION WHEN :

① X' IS SMOOTH ;

② $(X', \tilde{B} + \text{Exc}(\pi))$ IS SNC.

THEOREM [HIRONAKA] LOG RESOLUTION EXIST IN CHAR 0.

MINIMAL MODEL PROGRAM: A GENTLE INTRODUCTION

DEFINITION LET (X, B) BE A PAIR. A BIRAT'L MORPHISM

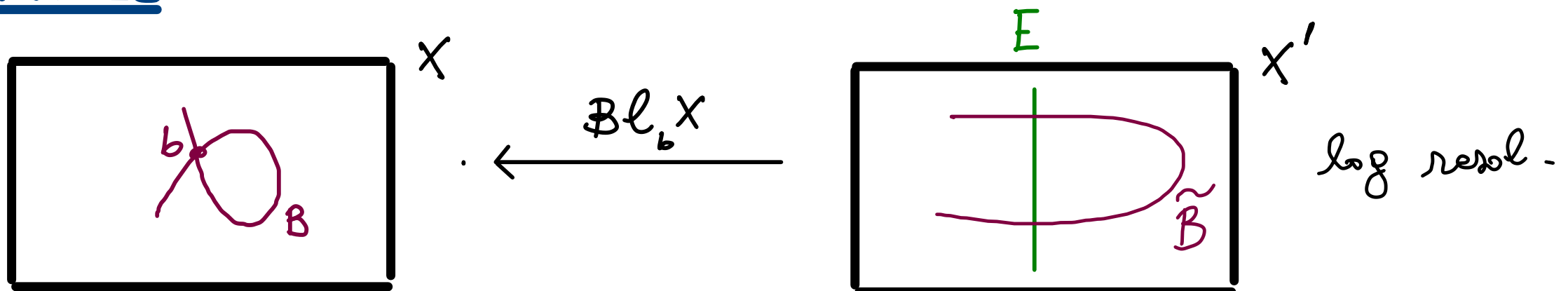
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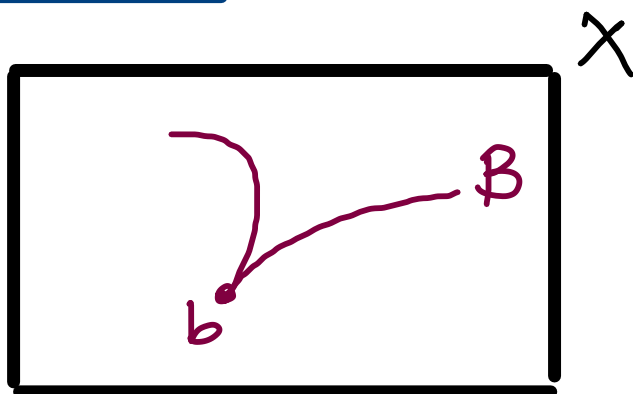
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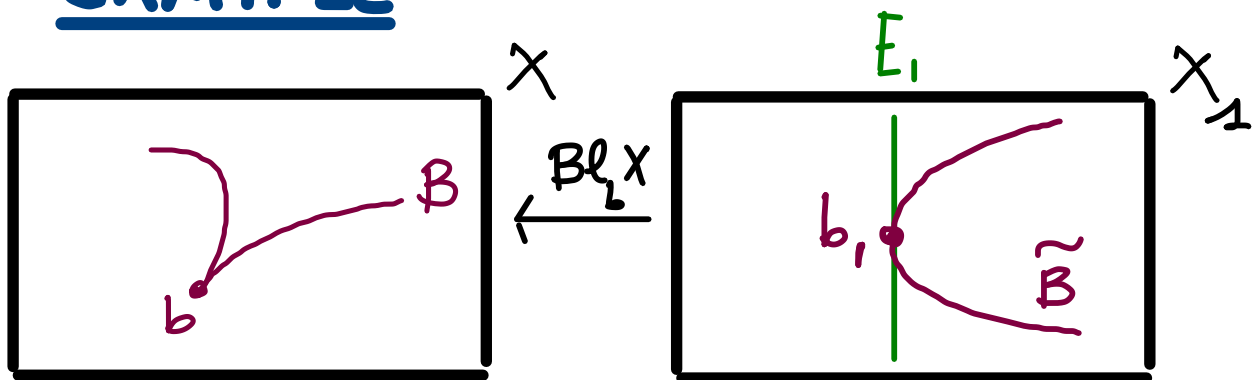
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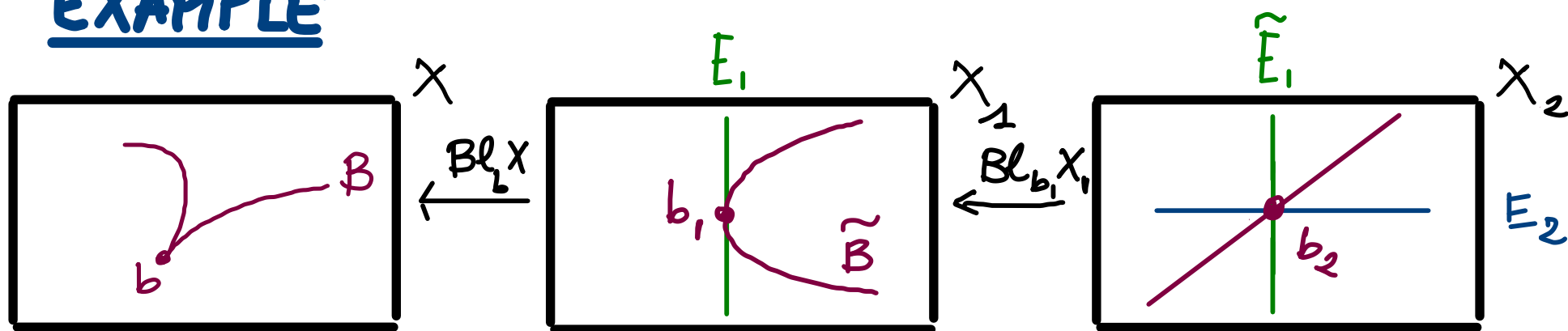
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EXAMPLE



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DEFINITION LET (X, B) BE A PAIR. A BIRAT'L MORPHISM

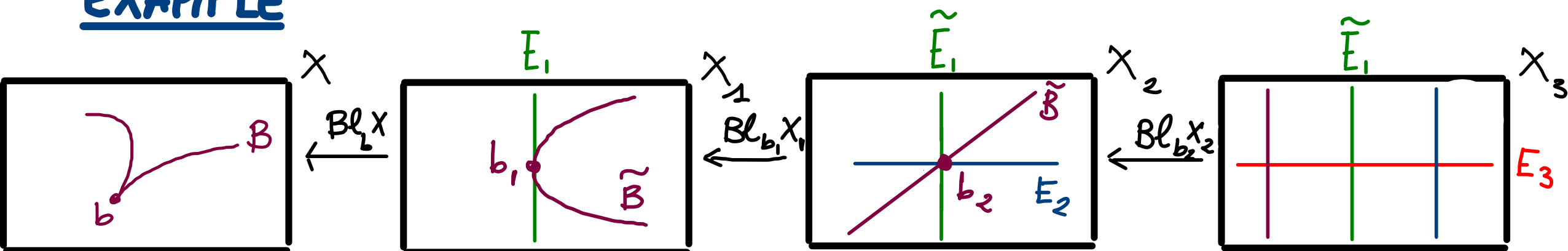
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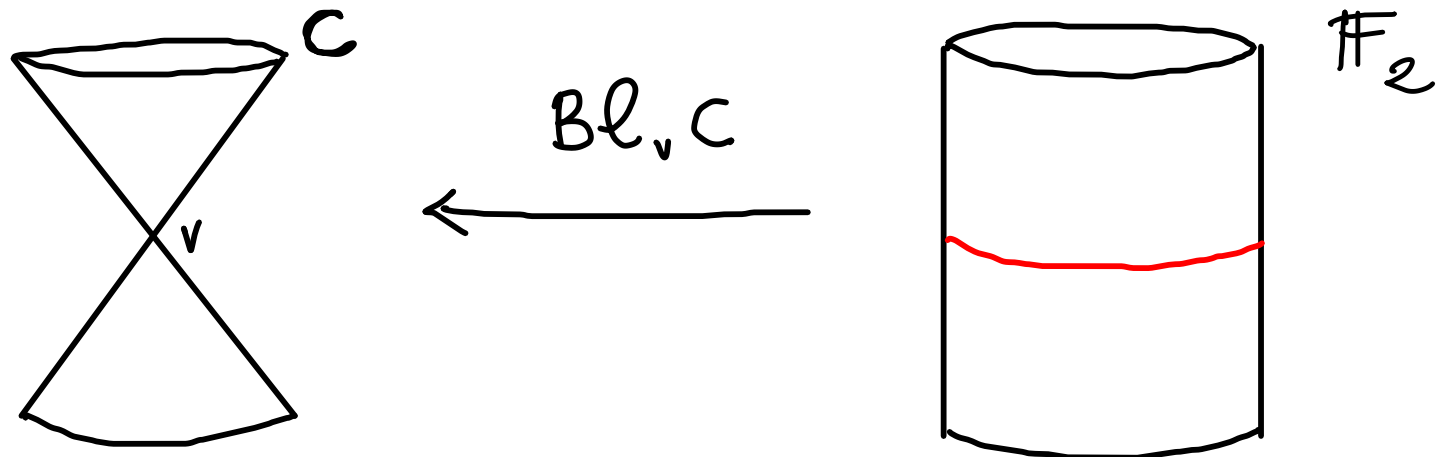
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WE SAW THAT WE WILL NEED TO CONSIDER SINGULAR OBJECTS.

TO DO THAT, WE NEED TO MEASURE THEIR SINGULARITIES.

WE WILL DO SO USING (BIRATIONAL) RESOLUTIONS.

GIVEN (X, B) A LOG PAIR AND A LOG RESOL.

$$\pi: X' \longrightarrow X$$

WE CAN ALWAYS WRITE IN A UNIQUE WAY

$$K_{X'} + \tilde{B} = \pi^*(K_X + B) + \sum_{\substack{\text{exc} \\ \text{div.}}} a_i E_i$$

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WE CAN ALWAYS WRITE IN A UNIQUE WAY

$$K_{X'} + B' = \pi^*(K_X + B)$$

$$\text{WHERE } B' = \pi_*^{-1} B + \sum_{\substack{\text{exc.} \\ \text{div.}}} a_i E_i$$

DEFINITION LET $D \subseteq X'$ BE A PRIME DIVISOR, THEN

$$a(D; X, B) := 1 - \text{coeff. of } D \text{ in } B'$$

log discrepancy

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

GIVEN (X, B) A LOG PAIR AND A LOG RESOL.

$$\pi: X' \longrightarrow X$$

WE CAN ALWAYS WRITE IN A UNIQUE WAY

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DEFINITION ① LET $D \subseteq X'$ BE A PRIME DIVISOR, THEN

$$a(D; X, B) := 1 - \text{coeff. of } D \text{ in } B'.$$

[THIS IS CALLED THE LOG DISCREPANCY OF D WRT (X, B)]

② THE TOTAL LOG DISCREPANCY OF (X, B) IS

$$a(X, B) := \inf a(D; X, B)$$

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

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② THE TOTAL LOG DISCREPANCY OF (X, B) IS

$$a(X, B) := \inf a(D; X, B)$$

SPECIAL DICHOTOMY : $a(X, B) = -\infty$

OR

$a(X, B) \geq 0$ MMP ☺

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

DEFINITION LET (X, B) BE A LOG PAIR.

THEN (X, B) IS

LOG CANONICAL

KLT KAWAMATA
LOG TERM.

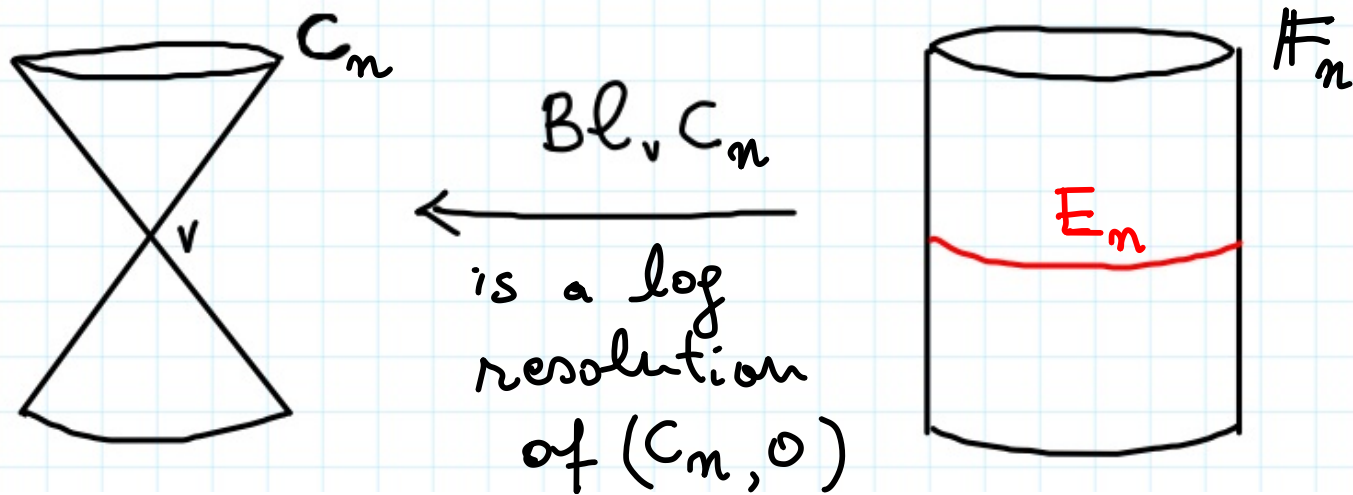
CANONICAL

TERMINAL

IF $q(X, B)$

≥ 0
 > 0
 ≥ 1
 > 1

EXAMPLE



$C_n =$ CONE OVER RAT'L NORMAL CURVE OF deg n

LET'S COMPUTE THE LOG DISCREPANCY OF E_n

$$K_{F_n} = \pi^*(K_{C_n}) + a_n E_n$$

$\left\{ \begin{array}{l} + E_n \text{ on both sides} \\ \end{array} \right. = a(E_n; C_n, 0)$

$$K_{F_n} + E_n = \pi^*(K_{C_n}) + (1 - a_n) E_n$$

$\left\{ \begin{array}{l} \cdot E_n \text{ on both sides} \\ \end{array} \right.$

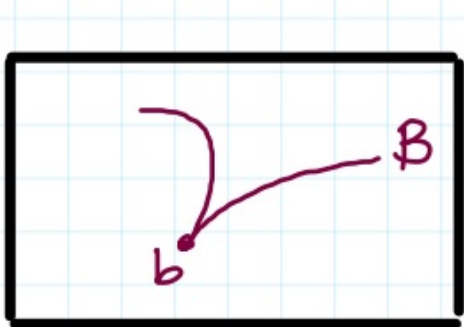
$$\frac{2}{n} = a(E_n, C_n, 0)$$

\uparrow

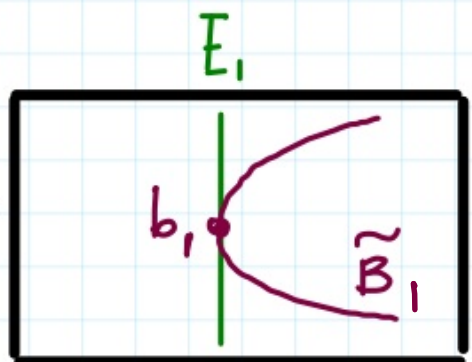
adj. to $E_n = \mathbb{P}^1$

$$-2 = (K_{F_n} + E_n) \cdot E_n = (\pi^*(K_{C_n}) + a(E_n, C_n, 0) E_n) \cdot E_n = 0 + a(E_n, C_n, 0)(-n)$$

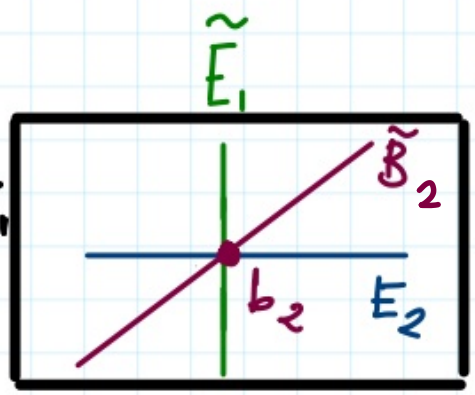
EXAMPLE



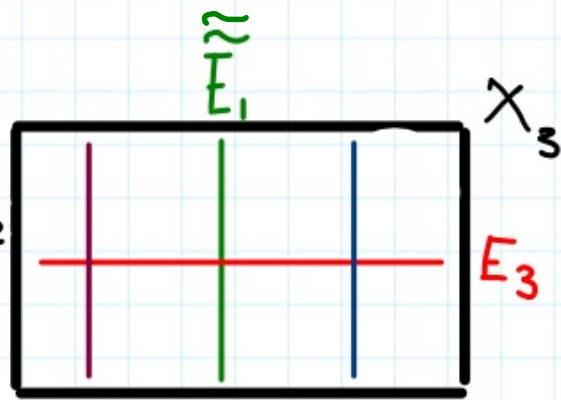
$$\begin{matrix} X \\ \leftarrow \frac{Bl_x}{\pi_1} \end{matrix}$$



$$\begin{matrix} X_1 \\ \leftarrow \frac{Bl_{b_1} X_1}{\pi_2} \end{matrix}$$



$$\begin{matrix} X_2 \\ \leftarrow \frac{Bl_{b_2} X_2}{\pi_3} \end{matrix}$$



$$\pi_1^* B = \tilde{B}_1 + 2E_1$$

$$\pi_1^* K_x = K_{x_1} - E_1$$

\Downarrow

$$\pi_1^* (K_x + B)$$

\parallel

$$K_{x_1} + \hat{B} + E_1$$

(X, B) IS NOT LC

$$\pi_2^* E_1 = E_2 + \tilde{E}_1$$

$$\pi_2^* \tilde{B}_1 = \tilde{B}_2 + E_2$$

$$\pi_2^* K_{x_1} = K_{x_2} - E_2$$

\Downarrow

$$\pi_2^* (K_{x_1} + \hat{B} + E_1)$$

\parallel

$$K_{x_2} + \tilde{E}_1 + E_2 + \tilde{B}_2$$

$$\pi_3^* \tilde{E}_1 = \tilde{\tilde{E}}_1 + E_3$$

$$\pi_3^* \tilde{B}_2 = \tilde{B}_3 + E_3$$

$$\pi_3^* E_2 = \tilde{\tilde{E}}_2 + E_3$$

$$\pi_3^* K_{x_2} = K_{x_3} - E_3$$

\Downarrow

$$\pi_3^* (K_{x_2} + \tilde{E}_1 + E_2 + \tilde{B}_2)$$

\parallel

$$K_{x_3} + \tilde{\tilde{E}}_1 + \tilde{\tilde{E}}_2 + \tilde{B}_3 + 2E_3$$

MINIMAL MODEL PROGRAM: A GENTLE INTRODUCTION

~~X SMOOTH PROJECTIVE VARIETY / \mathbb{C}~~

(X, B) LOG CANONICAL PAIR

CONJECTURE

STARTING FROM (X, B) , THERE EXISTS AN ALGORITHMIC WAY OF CONSTRUCTING A DIAGRAM OF THE FOLLOWING FORM

X — composition of special birat'l maps $\rightarrow (X', B')$



SUCH THAT f SATISFIES 1 OF THE FOLLOWING:

① f BIRAT'L & ~~K_X~~ is AMPLE

② f FIBRATION & $K_{X'} + B' = f^* H$ ample

③ f FIBRATION & $(K_{X'} + B')$ AMPLE on fibers

MINIMAL MODEL PROGRAM : A GENTLE INTRODUCTION

IN ORDER TO ACHIEVE SUCH RESULT, WE WILL TRY TO CONTROL & TUNE THE POSITIVITY/NEGATIVITY OF LOG DIVISORS.

(X, B) LOG PAIR

$K_X + B$ is \mathbb{R} -Cartier

$\forall C \subseteq X$ proper curve

$(K_X + B) \cdot C$

$K_X + B \cdot C \geq 0$

$\forall C \subseteq X$ proper curve

GOAL: either make $K_X + B$ nef after many birational transformations

OR

die trying (but in finite time)

THE CONE THEOREM

THEOREM [Mori, Kollár, Shokurov, ...]

LET (X, B) BE A LOG CANONICAL PAIR.

WE HAVE THE FOLLOWING DECOMPOSITION

$$\overline{NE}(X) = \overline{NE}(X)_{K_X + B \geq 0} + \overline{NE}(X)_{K_X + B < 0}$$

the closure the one of
effective
curves

THE CONE THEOREM

THEOREM

LET (X, B) BE A LOG CANONICAL PAIR.

WE HAVE THE FOLLOWING DECOMPOSITION

$$\overline{NE}(X) = \overline{NE}(X)_{K_X+B \geq 0} + \overline{NE}(X)_{K_X+B < 0} =$$

$$= \overline{NE}(X)_{K_X+B \geq 0} + \sum_{i \in I} \mathbb{R} + [C_i] \begin{array}{l} \text{extremal} \\ \text{rays} \end{array}$$

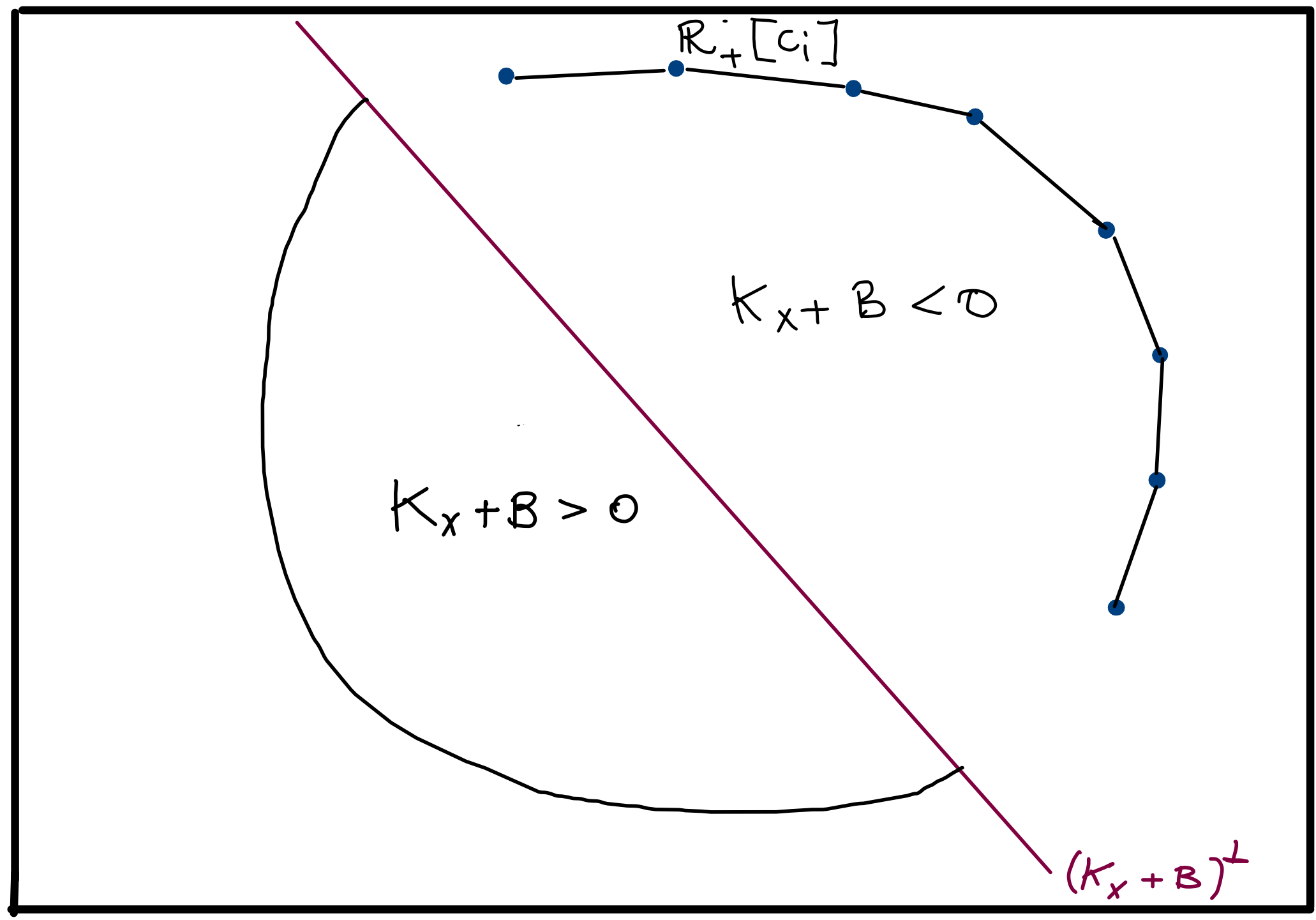
countable set (pointing to $i \in I$)
rat'l curves (pointing to $[C_i]$)
locally discrete away from $(K_X+B)^+$ (pointing to the sum)

$$-2 \dim X \leq (K_X+B) \cdot C_i < 0$$

HORIZONTAL SLICE OF $\overline{NE}(x)$

extremal rays $\stackrel{\text{def}}{=} v_1 + v_2 \in R_i$

\Downarrow
 $v_1, v_2 \in R_i$



THE CONE THEOREM

THEOREM

LET (X, B) BE A LOG CANONICAL PAIR.

WE HAVE THE FOLLOWING DECOMPOSITION

$$\begin{aligned} \overline{NE}(X) &= \overline{NE}(X)_{K_X+B \geq 0} + \overline{NE}(X)_{K_X+B < 0} = \\ &= \overline{NE}(X)_{K_X+B \geq 0} + \sum_{i \in I} \overbrace{R_i + [C_i]}^{R_i \text{ extremal ray}} \end{aligned}$$

\uparrow at most countable set \uparrow nat'l curves

$-2 \dim X \leq (K_X + B) \cdot C_i < 0$

MOREOVER, $\forall i \in I$, $\exists \text{cont}_{R_i} : X \longrightarrow \mathbb{Z}$ SUCH THAT

cont_{R_i} HAS CONNECTED FIBERS & IF $\text{cont}_{R_i}(c) = \text{pt} \Rightarrow [c] \in R_i$.